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# **Communications Subsystem DFS/NFS Comparison for the ECS Project**

**Technical Paper**

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formal review or government approval.**

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# Abstract

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The Communications Subsystem (CSS) provides the overall communications infrastructure, and the communications services to support other subsystems in the Science and Communications Development Office (SCDO) and the Flight Operations Segment (FOS). This document describes a comparison of DFS/NFS.

**Keywords:** CSMS, CSS, Communications, DCE, OODCE, Release A, DFS, NFS.

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## Abbreviations and Acronyms

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# 1. Introduction

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## 1.1 Background

With the transition from mainframe computers to networked machines comes a desire to be able to share files between the networked machines. Early methods of sharing files involved sending entire files between machines. NFS, developed by Sun, provided an alternative approach. Rather than transferring a file from one local disk to another, NFS allowed a user to transparently access files on a remote machine. Thus, when a user accessed a part of a file, his workstation would request that part of the file from the remote machine. This allowed a change to occur in computing. Diskless workstations were invented, they are able to use another machine's disk as though it were its own, however since disk prices have dramatically dropped this is not much of an issue any more. More importantly though, it allowed users to share data even when they were operating on different machines.

However, there were limitations to what the technology could do. Currently shared disk space needs have outgrown the original NFS design. Unfortunately, the current NFS designs still maintain backwards compatibility with the original design leaving several critical flaws in the technology. Many of these are addressed by AFS (which once stood for Andrew File System, but now is no longer an acronym for anything) and DFS (which stands for Distributed File System.)

The players:

NFS was developed by Sun, and is currently the most widely used UNIX networked file system. While it has had some improvements over the past decade, the technology is still basically the same as when it was first developed. While it allows transparent sharing of files, it provides only host based security effectively - thus the only real restrictions you can place on who can read or modify files is based on what machine a user is on. While it currently provides some caching of files for faster access, it can not distinguish between part of a file being changed or the entire file being changed. For example if a user reads bytes 1-100 of a file, and her coworker modifies bytes 1024-2048, if she re-reads bytes 1-100 those bytes will have to be sent across the network again. Also, each time a user accesses a file the system either checks with the server to see if the file has changed (causing network delays) or risks the possibility that the file has changed since it was last checked. Usually a time period of 10-30 seconds is maintained, inside of which it is assumed that no files have changed. Another limitation of NFS is that each client workstation must maintain its own 'image' of how networked files are accessed, and while it is possible administratively to keep each client machine in sync with the others to provide the same view of the file system, it is not automatic. Also, if a file server is moved, each client must be re-configured in order to be able to find the new server.

AFS was developed by jointly by IBM and Carnegie Mellon University. It later became its own company called Transarc. Transarc was subsequently entirely bought by IBM and is now a fully owned subsidiary. AFS provided real security based on Kerberos authentication. It allows ACLs (Access Control Lists) to be placed on each directory specifying which users are given what types

of access to the files in that directory. It also allows for AFS users in other Kerberos cells (very similar to DCE cells) to authentically access files in other cells. AFS heavily relies on caching of files for performance. When a client receives a part of a file, it saves it either in disk or memory, so that subsequent accesses to that file can be handled locally instead of requiring network traffic. It also establishes a guarantee with the file server that if the file changes, the client will be notified. This prevents the client from having to constantly poll the file server to see if there have been changes. However, when there is a change to the file, AFS must re-fetch any needed parts since it also doesn't understand about changes to one part of a file not effecting other parts of the file. In stark contrast to NFS, AFS maintains a global image of the filesystem. This is maintained on the file server machines, and is automatically shared by all AFS clients. Thus, when re-configuring, moving, or adding file servers, no changes need to be made to any of the client machines, and all client machines will immediately be able to access the servers. One feature that AFS includes which also can dramatically improve performance is the idea of replication. Multiple file servers can store the same sets of data (if the data is read only), and clients will randomly select from the available servers before requesting files. This provides load balancing, improved performance and fault tolerance all at once. So, if one of the replicated file servers is down, the user's computer will automatically contact one of the other servers that is still up and continue work without interruption or failure.

DFS was designed by the OSF with contributions from Transarc. DFS is very similar to AFS, and here I will only list the differences between DFS and AFS. DFS uses DCE's security (which also happens to be based on Kerberos.) Unlike AFS, DFS provides ACLs for each file, rather than for each directory. However, currently it has the limitation that only users inside of the file server's DCE cell can have authenticated access to the files. DFS provides partial caching as well, but it is able to determine when modifications to other parts of a file do not affect currently used parts of the file and can prevent un-needed network data transmission. Currently there are other implementations of DFS (since it is an OSF standard) but Transarc/IBM are the only ones-at the time of this writing-to support ACLs and replication. However, other vendors are planning to support it in their servers in the future.



## 2. Objectives

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### 2.1 Objectives

A remote file system allows remote users to access files on a disk as though the files were locally available. This provides a transparent method to transfer files when needed between machines. Also, when compared to FTP, it allows partial files to be transferred rather than requiring the entire file to be transferred.

This analysis studies features and evaluates performance of two remote file systems: DFS and NFS (V2).

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## **3. Requirements and Assumptions**

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### **3.1 Requirements**

There is currently a requirement to use DFS as a pull area for Release A.

### **3.2 Assumptions**

DFS and AFS require proprietary software on the clients to be able to access files, while NFS clients are standard on just about every know UNIX workstation.

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## 4. Description of Alternatives

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### 4.1 Description

NFS (V2) servers across platforms and vendors provide the same functionality. DFS servers however offer different functions based on the vendor of the software. This evaluation will focus on Transarc's release of the DFS server for the Sparc family. Thus, for this study, the two setups looked at are:

1. An NFS based file server.
2. A DFS based file server.

This trade study initially evaluates the two alternatives as the basic underlying configuration and then evaluates the benefits gained by Transarc's DFS servers over other vendor's DFS servers, with mention of AFS. AFS is only being mentioned in passing because of the existing DCE cell architecture, it makes much more sense going with DFS over AFS.

### 4.2 NFS Based File Server

In this configuration there would be a single NFS server which could provide files to clients. Clients would have to explicitly mount the file server before being able to access the files. NFS clients are available for nearly all UNIX workstations and can optionally provide some caching of files on the client side. Most vendors have NFS servers available for their systems.

### 4.3 DFS Based File Server

In this configuration, there would be one or more DFS servers which could provide files to clients. Clients would need to be running on a machine with the DCE client software installed. File caching is automatically done. If desired users could be required to authenticate themselves before gaining access to files.

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## 5. Evaluation

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### 5.1 Description

In evaluating the alternative solutions for the cell configuration, the following criteria were applied.

Cost

Performance

Security

Scalability

The weights in the table below are assigned as follows:

- 1) relatively minor
- 2) significant
- 3) critical

The score given to a server is assigned as follows:

- 0) Extremely poor
- 1) Poor
- 2) Average
- 3) Good

**Table 5-1. Cell Configuration Evaluation Matrix**

	Weight	NFS	DFS
<b>Cost</b>			
<b>Software</b>	2	2	1
<b>Ease of administration</b>	1	1	3
<b>Performance</b>			
<b>Response time</b>	3	3	2
<b>Reliability of connections</b>	2	1	2
<b>Scalability</b>			
<b>Additional replicated severs</b>	1	0	3
<b>Security</b>			
<b>Security</b>	3	1	3
<b>Authentication across cells</b>	3	1	2
<b>Total</b>		<b>22</b>	<b>33</b>

In this evaluation, the following criteria were applied:

#### Cost

Usually NFS servers are shipped with the Vendor OS and as such, are free. DFS servers cost approximately \$2,000 per server. NFS clients are usually free, and DFS clients are distributed as part of the DCE client. Thus, many platform support it (including, among others, DEC-OSF/1, HP-UX, AIX, Solaris, Cray, and this list is expanding.)

#### Operational cost

DFS is substantially easier to administer than NFS. Also, DFS allows delegation of certain aspects of the file system to some people while still maintaining control over the rest of the system. With NFS there is very little differentiation between different types of administration.

#### Performance

NFS performance is roughly twice the speed of DFS performance on a local area network under small file system load. (4 clients to 1 server) It is expected, although untested, that as the load increases DFS performance will improve and surpass the performance of NFS. DFS also allows replication of data, so that two or more servers may contain the same data set. This would allow the clients to spread the load between servers.

#### Scalability

Addition or renaming of NFS servers may not be done transparently. Clients must know the name of each machine that is a file server before being able to request any files from that machine. With DFS, addition of machines (and replicas of data) may be done transparently, and without any downtime for the clients.

#### Security

NFS, while in theory, allows group access controls to files, in practices is fundamentally insecure because it is based on an insecure protocol. DFS provides Kerberos authentication integrated with DCE. Access per file and directory may be restricted based on user, and may be administered separately from the basic DFS administration if desired.

## 5.2 DFS Variations

There are alternative vendors for DFS servers for different operating systems. Currently most of them do not support ACLs or replication. Most of the vendors have plans to eventually support these features, but not for a few years, which is why this study concentrated on Transarc's implementation.

Transarc also sells AFS servers and clients. AFS is similar to DFS, and in many ways was the parent of DFS. The main differences are that AFS uses vanilla Kerberos for security, rather than DCE's Kerberos, and that it only provides ACLs per directory rather than per file.



## 6. Conclusions / Recommendations

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### 6.1 Conclusions and Recommendations

Based on the above evaluation, DFS is recommended over NFS for any networked file system uses. However, it is not clear if DFS is currently mature enough to trust for critical applications. It will probably be quite stable 1-2 years from now.

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# Abbreviations and Acronyms

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ACL	Access Control List
AFS	Andrews File System
AI&T	Algorithm Integration and Test
AIT	Algorithm Integration Team
ANSI	American National Standards Institute
API	Application program (or programming) interface
ASCII	American Standard Code for Information Exchange
ATM	Asynchronous Transfer Mode
ARP	Address Resolution Protocol
ASTER	Advanced Spaceborne Thermal Emission and Reflection Radiometer
BB	Bulletin Board
BBS	Bulletin Board Service
BIND	Berkeley Internet Name Domain
BGP	Border Gateway Protocol
BOA	Basic Object Adapter
CAC	Command and Activity Controller
CCB	Change Control Board (Hughes Convention)
CCB	Configuration Control Board (NASA Convention)
CCR	Configuration Change Request
CDS	Cell Directory Service
CDR	Critical Design Review
CDRL	Contract data requirements list
CERES	Clouds and Earth's Radiant Energy System
CIDR	Classless Interdomain Routing
CM	Configuration management
CMAS	Configuration Management Application Service
CMIP	Common Management Information Protocol
CNE	Campus Network Environment

CORBA	Common object request broker architecture
COTS	Commercial off-the-shelf (hardware or software)
CPU	Central processing unit
CSMS	Communications and System Management Subsystem
CSS	Communication Subsystem
DAAC	Distributed Active Archive Center
DADS	Data Archive and Distribution System
DB	Database
DBMS	Database management system
DCE	Distributed computing environment (OSF)
DEC	Digital Equipment Corporation
DECOM	FOS Decommutation Process
DFS	Distributed File System
DID	Data item description
DME	Distributed Management Environment
DNS	Directory Name Service
DOF	Distributed Object Framework
DPR	December Progress Review
DS	Data Server (FOS)
DTS	Distributed Time Server (part of DCE)
ECS	EOSDIS Core System
EDOS	EOS Data and Operations Center
EDF	ECS Development Facility
E-Mail	Electronic Mail
EMC	Enterprise Monitoring and Coordination
EOC	EOS Operations Center (ECS)
EOS	Earth Observing System
EOSDIS	Earth Observing System Data and Information System
EP	Evaluation Prototype
ESN	EOSDIS Science Network
EPV	Endpoint Vector

FDDI	Fiber distributed data interface
FDF	Flight Dynamics Facility
FOS	Flight Operations Segment
Ftp	File Transfer Protocol
GB	Gigabyte ( $10^9$ )
GCDIS	GDS Global Directory Service
GDS	Global Directory Service
GSFC	Goddard Space Flight Center
GUI	Graphic user interface
HAIS	Hughes Applied Information Systems (ECS)
HiPPI	High Performance Parallel Interface
HP	Hewlett Packard
Http	Hyper Text Transfer Protocol
I/F	Interface
I&T	Integration & Test
IBM	International Business Machines, Inc.
ICD	Interface control document
ICMP	Internet Control Messaging Protocol
IDL	Interface Definition Language
IEEE	Institute of Electrical and Electronics Engineers
IETF	Internet Engineering Task Force
IP	Internet Protocol
IR-1	Interim Release 1
ISO	International Standards Organization
ISO+	IsoCELL (Isolation Cell)
ISS	Internetworking Subsystem of CSMS
IST	Instrument Support Toolkit
IST	Instrument Support Terminal
Kerberos	Security protocol developed by MIT; base for DCE security
Kftp	Kerberized file transfer protocol
KLOC	Kilolines ( $10^3$ ) of code

Ktelnet	Kerberized telnet
LAN	Local area network
LaRC	Langley Research Center
LLC	Logical Link Control
LOC	Lines of code
LSM	Local System Management
M&O	Maintenance and operations
MBONE	Multicast Backbone
MIB	Management Information Base
MIME	Multimedia Internet Mail
MLM	Mid-Level Manager
MOPITT	Measurement of pollution in the troposphere
MOSPF	Multicast Open Shortest Path First
MR-AFS	Multi-Resident Andrew File System
MSFC	Marshall Space Flight Center
MSS	Systems Management Subsystem
MUI	Management User Interface
NCR	Non-conformance Report
NFS	Network file system
NIC	Network Interface Card
NNTP	Network New Transfer Protocol
NOAA	National Oceanic and Atmospheric Administration
NOLAN	Nascom Operational Local Area Network
NSI	NASA Science Internet
NTP	Network Time Protocol
OA	Off-Line Analysis Process
OLAP	On-Line Analytical Processing
OLTP	On-Line Transaction Processing
OMG	Object Management Group
OMT	Object Modelling Technique
OO	Object-oriented

OODCE	Object-oriented DCE
OODBMS	Object-oriented database management system
ORB	Object Request Broker
OS	Object Services (CSS building blocks)
OSF	Open Software Foundation
OSI	Open System Interconnect
OSI-RM	OSI Reference Model
OSPF	Open Shortest Path First
PAC	Privilege Attribute Certificate
PDR	Preliminary Design Review PDR-A
PDU	Protocol Data Unit
PPP	Point-to-Point Protocol
POSIX	Portable Operating System Interface for Computer Environments
PSC	Pittsburgh Supercomputing Center
PTGT	Privilege Ticket Granting Ticket
RDBMS	Relational database management system
RFA	Remote File Access
RFC	Request for comments
RIP	Routing Information Protocol
RMA	Reliability, Maintainability, Availability
RMON	Remote Monitoring
RMP	Reliable Multicast Protocol
RPC	Remote procedure call
RTS	Real-Time Server (FOS)
SCF	Science Computing Facility
SDPF	Sensor Data Processing Facility
SDR	Software/System Design Review
SDR	Sensor data record
SGI	Silicon graphics
SLOC	Source lines of code
SMC	System Monitoring and Coordination

SMDS	Switched Multi-megabit Data Service
SMTP	Simple Mail Transfer Protocol
SNMP	Simple Network
SQL	Simple Query Language
TBD	To be determined
TCP/IP	Transmission Control Protocol/Internet Protocol
TGT	Ticket Granting Ticket
TMN	Telecommunications Management Network
TRMM	Tropical Rainfall Measurement Mission
TSDIS	TRMM Science Data Information System
UDP	User Datagram Protocol
UIOAR	User Interface Off-Line Analysis Request Window
URL	Universal Resource Locator
US	User Station (FOS)
UUID	Universal Unique Identifier
UTC	Universal time code
V0	Version 0
VT	Virtual Terminal
WAN	Wide area network
WWW	World Wide Web
X	X Protocol
X.500	OSI standard for directory services (207)
XDS	X/Open Directory Service
XFN	X/Open Federated Naming
XOM	X/Open OSI-Abstract-Data Manipulation